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The Constriction Decomposition Method for Coverage Path Planning

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Stanley Brown* and Steven L. Waslander†

Abstract—The task of coverage path planning in 2D indoor and outdoor environments is classified as a NP-hard problem, and has been an active research topic for over 30 years. We derive a novel, exact cellular decomposition method called the Constriction Decomposition Method (CDM) and apply it to complex indoor environments. The CDM rapidly identifies constriction points in the environment and decomposes the environment into easily traversable cells by exploiting the geometric information contained in the environment's straight skeleton. Several heuristic path planning methods that find paths that completely cover each cell are explored. We apply our method on a complex indoor office-like environment and compare our results to existing methods.

I. INTRODUCTION

Coverage path planning (CPP) is the task of determining an optimal set of paths that allow a robot or agent to transverse all free space in an environment. CPP has numerous applications for tasks that require the entirety of free space in an environment to be covered, such as crop harvesting, inspection, and mapping. The focus of this paper is the derivation of a CPP algorithm that produces coverage paths tailored to operation in complex human environments with corridors, rooms and islands, by identifying constriction points in known maps at which to divide the environment in a logical manner.

The CPP problem has been an active research topic since at least the 1980s [1] and several main strategies for approaching the problem have emerged. In many cases, the environment is broken down, either explicitly or implicitly, into cells or sub-regions to simplify the generation of coverage paths. If the cells are non-intersecting and the union of all cells fill the entire environment, the decomposition is called an exact cellular decomposition [2]. Examples of exact cellular decomposition include the Boustrophedon decomposition [3], Morse cell decomposition [2], minimum sum of angles (MSA) decomposition [4] or greedy convex polygon decomposition [5]. The Spiral Spanning Tree Coverage (Spiral-STC) method by [6] is an example of an approximate cellular decomposition, which is also simultaneously classified as a grid based method.

CPP algorithms can be classified as either offline or online methods [1]. Online CPP algorithms assume that limited or no prior knowledge of the environment exists and usually rely on a set of heuristic methods to ensure complete coverage [1]. Examples online CPP algorithms can be found in works of [5] and [7].

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In this paper, we propose an offline method that generates a coverage path plan using a novel, exact cellular decomposition called the constriction decomposition method (CDM). Our method works by decomposing a pre-mapped environment into a set of cells based on constriction points, which are defined as areas where the narrowest point is less than the two or more neighbouring areas. For indoor environments, a hallway between two large rooms can be thought of as a constriction point and by decomposing the environment into cells based on these points, an intuitive, room-based decomposition results. It is also shown that the resulting cells can be completely covered using a set of simple contour following paths, followed by a series of back and forth paths that are aligned with the longest edge of the cell.

II. RELATED WORK

One of the earliest exact cellular decomposition methods developed is the trapezoidal decomposition method which breaks polygonal environment down into cells by inserting a set of edges up and down from each obstacle vertex to the bounding polygon [8]. Every cell that is generated in this fashion is trapezoidal, convex and can be covered in a series of back and forth motions that are parallel with the direction of the inserted edges. While the trapezoidal decomposition method is shown to be algorithmically complete, it typically generates a number of cells that could be merged and still be covered using the same coverage pattern. This idea is formalized in the derivation of the Boustrophedon Decomposition [3].

The word boustrophedon comes from ancient Greek and literally means "the way of the ox" [3]. In the Boustrophedon Decomposition, a set of cells are generated assuming that the robot covers a space in a series of back and forth motions similar to the way farmers plow or work a field. By running a scan line that is parallel to the direction of the back and forth motions the robot will take when covering an environment, and inserting an edge whenever an edge can be drawn both above and below the vertex of an obstacle, a set of cells is created. These cells can always be covered using a boustrophedon like coverage pattern and by visiting all the cells the method is shown to be complete [3].

In later work, Choset [2] demonstrated that the Boustrophedon Decomposition can be generalized by proposing a cellular decomposition based on the critical points of Morse functions. The Morse Decomposition can be generalized to any n-dimensional space [2] and also allows for different coverage patterns to be used such as spiral, circular, or diamond patterns. In the case of the Boustrophedon Decomposition, its coverage pattern can be represented by the

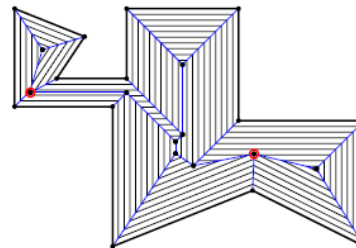


Fig. 1: An example straight skeleton in a non-convex polygon along with offset polygons located at $t = \{1, 2, \dots, 10\}$. Skeleton arcs are highlighted in with thin blue lines, skeleton nodes are denoted by black dots and nodes that correspond to split events are circled in red.

to move along the angular bisector of adjacent edges, forming edges in the straight skeleton. This process continues as long as the boundary does not change topologically.

During the shrinking process, there are two possible topological changes that can occur: an edge event and a split event [13]. An edge event occurs when an edge, e_i , shrinks to a length of zero, which causes its neighbouring edges to become adjacent. A split event occurs when a reflex vertex of the wavefront collides with an edge and causes it to split into two new edges, which in turn split the wavefront into two separate polygons. The shrinking process then still continues in all polygons until all edges shrink to a length of zero. The bisector line segments traced out by the vertices of the edges during the entire process are called the skeleton edges, and the bisector arc start and end points are called the nodes of $S(F)$. A node is said to be a contour node if it corresponds to a vertex in V and a skeleton node if it does not.

The shrinking process also gives rise to a hierarchy of nested polygons, a subset of which are shown in Figure 1 for an example polygonal environment along with the skeleton arcs, edge and split nodes. In Figure 1, it can be noted that the location of a split node gives rise to an additional polygon at the corresponding time steps. For the example polygon, the first split occurs at $t = 3.6$ and second split occurs at $t = 5.8$.

The process of shrinking a polygon for some time t at speed, ω , can be related to the task of creating a set of contour following paths. If an agent with a width, w , performing a coverage operation was to follow the contours of the polygonal environment m times, the polygon introduced for the m^{th} pass is the same as the polygon created by the shrinking process detailed above at a time $t = w \cdot m$, and indeed, $\omega = w$ in this case. Therefore, the coverage problem time that each node in $S(F)$ was created corresponds to the

number of times an agent can circle the free space of an environment without overlapping previous paths.

The location and time where the split nodes are created are of the greatest interest to our work as they encode the constriction points of an environment. Whenever a split event occurs, the resulting polygon of the wavefront is split into two separate polygons which means that the nodes on either side of the split node are created at a later time. Therefore, if the polygon is decomposed based on the split nodes in its straight skeleton, it is possible to create a set of cells that have no split nodes in their straight skeletons. That is, for such polygons, it will always be possible generate contour following paths that spiral inwards without crossing over a previous path, as described below. We term a decomposition based on this method the Constriction Decomposition Method (CDM).

B. The Constriction Decomposition (CDM)

Let N denote the set of all nodes, n_i , in the straight skeleton of the free space, $S(F)$, such that $n_i \in N$. Each node, n_i , is defined by its location in F , and has an associated time of creation, t_n , node type, $\tau_i = \{\tau_i^k, \tau_i^s\}$, (either skeleton or contour), event type, $\nu_i = \{\nu_i^s, \nu_i^e\}$, (split or edge) and a set of neighbouring nodes, M_i , connected by an edge in the straight skeleton graph. Given any node in the graph, its neighbours and their corresponding properties can be determined using graph traversal methods.

One method to find a set of decomposition edges E' such that F is decomposed into a set of cells C based on N is to simply search the skeleton of F , $S(F)$, for any split nodes. When a split node, n_{split} , is found, a new vertex, v_{new} , is inserted into the edge $e_{n,m}$ impacted by the angular bisector of v_i such that the length of e_{new} is minimized. In the case where the nearest point on the edge is equal to one of its vertices, e_{new} , is inserted between that vertex and v_i and no additional vertex is added to F . The insertion of the edge e_{new} results in a new cell $c_i \in F$ being created and also removes a split event from the skeleton of the remaining, non-decomposed part of $F' = F \setminus c_i$. This process is repeated until no split nodes exist and F is completely decomposed into a set of cells C . The CDM process is summarized algorithmically below and applied to the polygon in Figure 2.

Let $\theta : N \rightarrow \{\tau_i^k, \tau_i^s\}$ be function that takes a node $n_i \in N$ and returns its type, either contour or skeleton, and let $\psi : N \rightarrow \{\nu_i^s, \nu_i^e\}$ be a function that takes a node and returns its event type. Defining $\eta(v_i, e_{n,m})$ to be a function that takes a vertex, $v_i \in V$, and an edge, $e_{n,m} \in E$, and returns a virtual vertex, v^* on the edge, $e_{n,m}$ such that the minimum distance between v_i and v^* is minimized. Then the CDM process can be described as follows in Algorithm 1.

The Constriction Decomposition Method for Coverage Path Planning

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Abstract—The task of coverage path planning in 2D indoor and outdoor environments is classified as a NP-hard problem, and has been an active research topic for over 30 years. We derive a novel, exact cellular decomposition method called the Constriction Decomposition Method (CDM) and apply it to complex indoor environments. The CDM rapidly identifies constriction points in the environment and decomposes the environment into easily traversable cells by exploiting the geometric information contained in the environment's straight skeleton. Several heuristic path planning methods that find paths that cover each cell are explored. We apply our method of decomposition to a set of test environments and compare our results to other methods.

I. INTRODUCTION

Coverage path planning (CPP) is the task of determining an optimal set of paths that allow a robot or agent to traverse the free space in an environment. CPP has numerous applications including autonomous crop harvesting, industrial inspection, and pipeline inspection. The derivation of a CPP algorithm that produces coverage paths tailored to operation in complex human environments with corridors, rooms and islands, by identifying constriction points in known maps at which to divide the environment in a logical manner is a long standing problem.

The CPP problem has been a major research focus since at least 1980 [1]. In the early days of research, the approach to the problem were heuristic. In many cases, the environment is broken down, either explicitly or implicitly, into cells or sub-regions to simplify the generation of coverage paths. If the cells are non-intersecting and the union of all cells fill the entire environment, the decomposition is called an exact cellular decomposition. The CDM is an exact cellular decomposition in the sense that the union of all cells covers the entire environment. The CDM is based on the Morse decomposition [3]. Morse decomposition [4] or greedy convex polygon decomposition [5]. The Spiral Spanning Tree Coverage (Spiral-STC) method by [6] is an example of an approximate cellular decomposition, which is also simultaneously classified as a grid based method.

CPP algorithms can be classified as either offline or online methods [1]. Online CPP algorithms assume that limited or no prior knowledge of the environment exists and usually rely on a set of heuristic methods to ensure complete coverage [1]. Examples online CPP algorithms can be found in works of [5] and [7].

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The word boustrophedon comes from ancient Greek and literally means "the way of the ox" [3]. In the Boustrophedon Decomposition, a robot moves in a zig-zag pattern, covering the entire environment. The Boustrophedon Decomposition is a novel method that decomposes the environment into cells based on constriction points. The Boustrophedon Decomposition is a novel method that decomposes the environment into cells based on constriction points. The Boustrophedon Decomposition is a novel method that decomposes the environment into cells based on constriction points.

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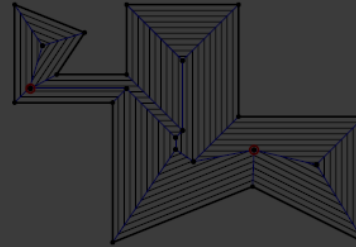


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II. RELATED WORK

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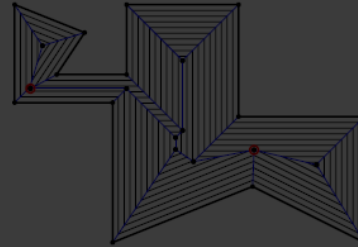


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coverage path planning



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[A Bircher](#), [M Kamel](#), [K Alexis](#), [M Burri](#) ... - *Autonomous ...*, 2016 - Springer

모든 날짜

Abstract This paper presents a new algorithm for three-dimensional coverage path planning for autonomous structure inspection operations using aerial robots. The proposed approach is capable of computing inspection paths in 3D environments (e.g., buildings) that are 17회 인용

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[E Galceran](#), [R Campos](#), [N Palmesras](#) ... *Journal of Field ...*, 2015 - Wiley Online Library

Abstract This paper presents a new algorithm for three-dimensional coverage path planning for autonomous structure inspection operations using aerial robots. The proposed approach is capable of computing inspection paths in 3D environments (e.g., buildings) that are used in a 2.5-dimensional (2.5 D) prior bathymetric map to plan a nominal coverage path that 18회 인용

dealing with the topic

“Coverage Path Planning on 2D Map”

관련도별 정렬

날짜별 정렬

Energy-aware **coverage path planning** of UAVs

[C Di Franco](#), [G Buttazzo](#) - *Autonomous Robot Systems and ...*, 2015 - [ieeexplore.ieee.org](#)

모든 언어

한국어 웹

Abstract: **Coverage path planning** is the operation of finding a **path** that covers all the points of a specific area. Thanks to the recent advances of hardware technology, Unmanned Aerial Vehicles (UAVs) are starting to be used for photogrammetric sensing of large areas in 21회 인용

특히 포함

서지정보 포함

[인용] A new 3D **Coverage Path Planning** Approach for Agricultural Robots to Minimize Skip Areas between Swaths over Fields of Complex Terrain
[I Hameed](#), [A la Cour-Harbo](#) - *Robotics and Autonomous Systems*, 2015

**2. Coverage path planning
with topological analysis
is quite interesting.**

I hope...

CPP implies Space Efficiency

Maybe no direct relation.
At least weak relation...!

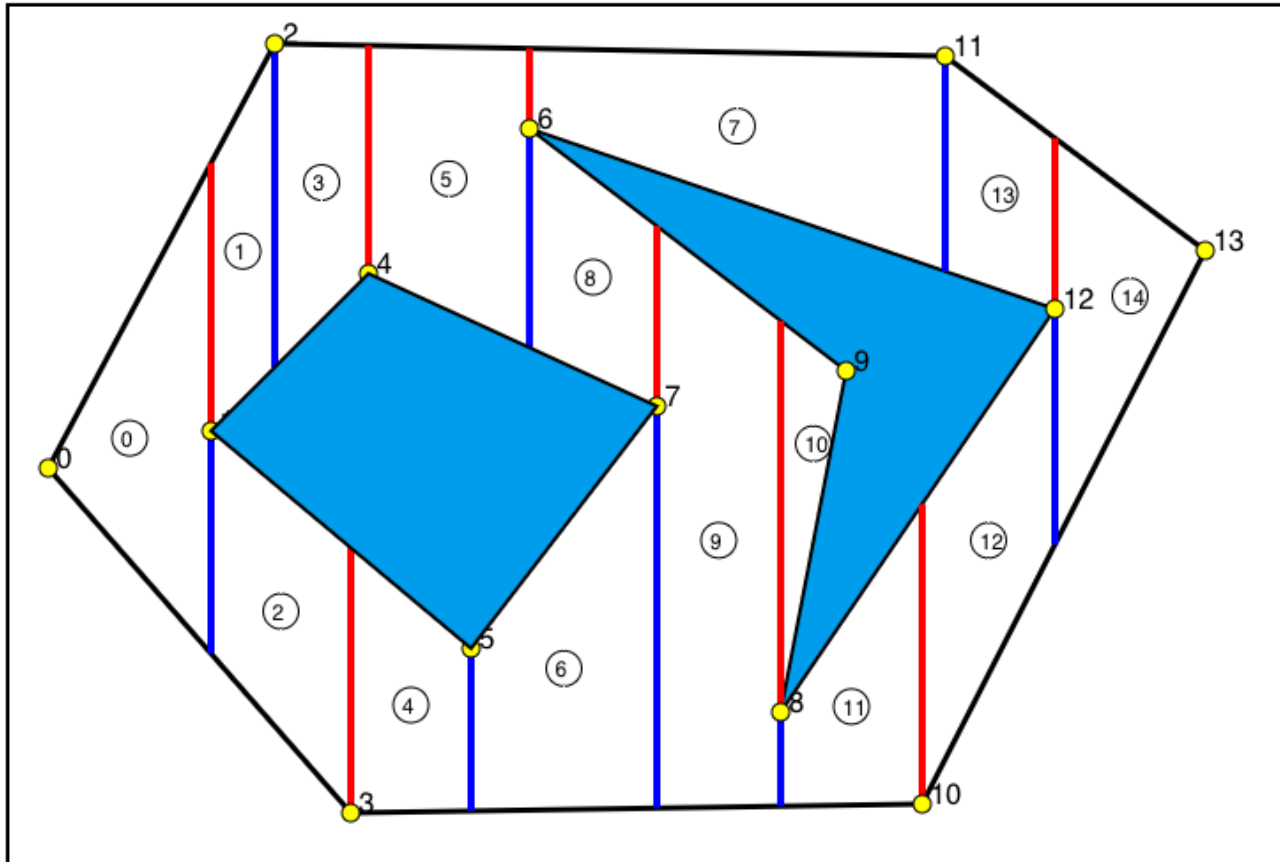
Some Critical Points Inside the Polygon Map



The **Constriction Decomposition** Method for **Coverage Path Planning**

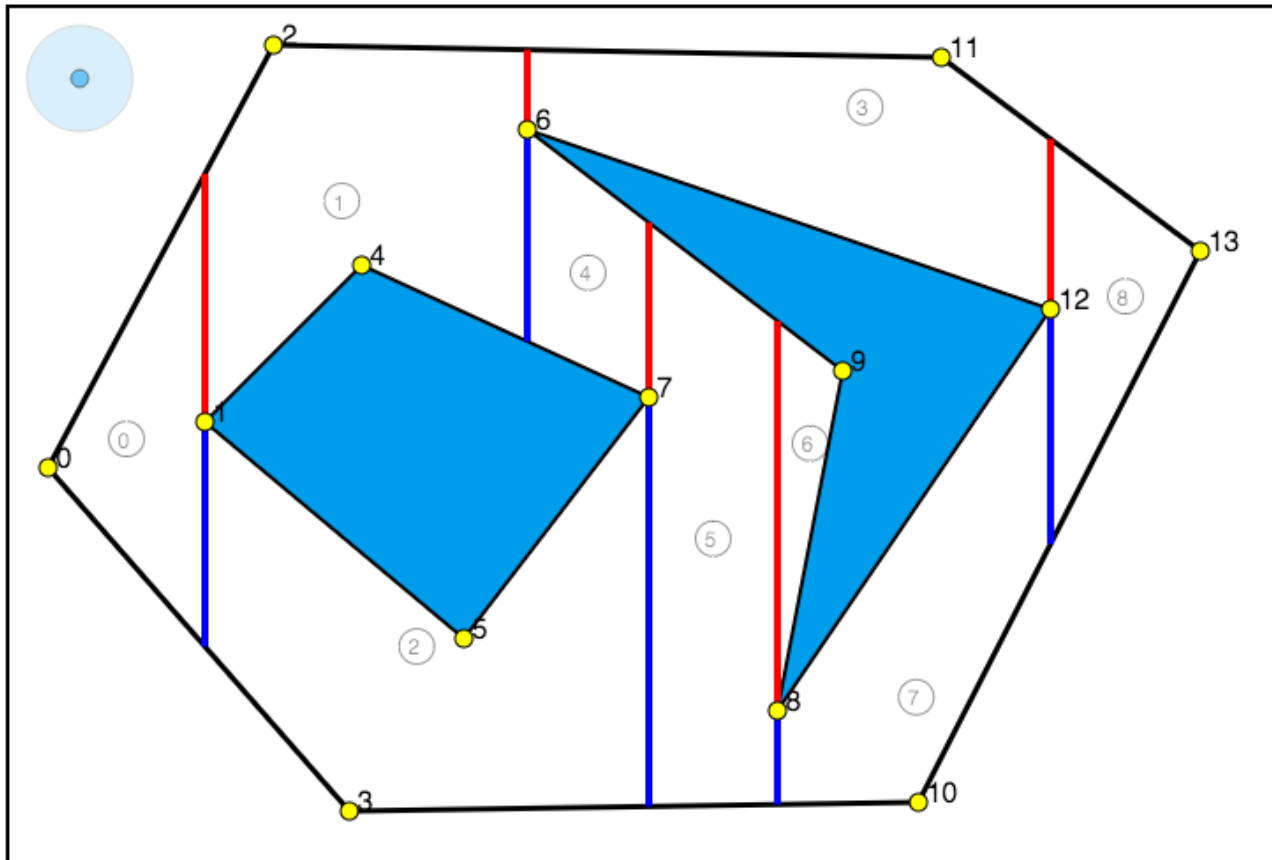
[Previous Work]

Trapezoidal Decomposition

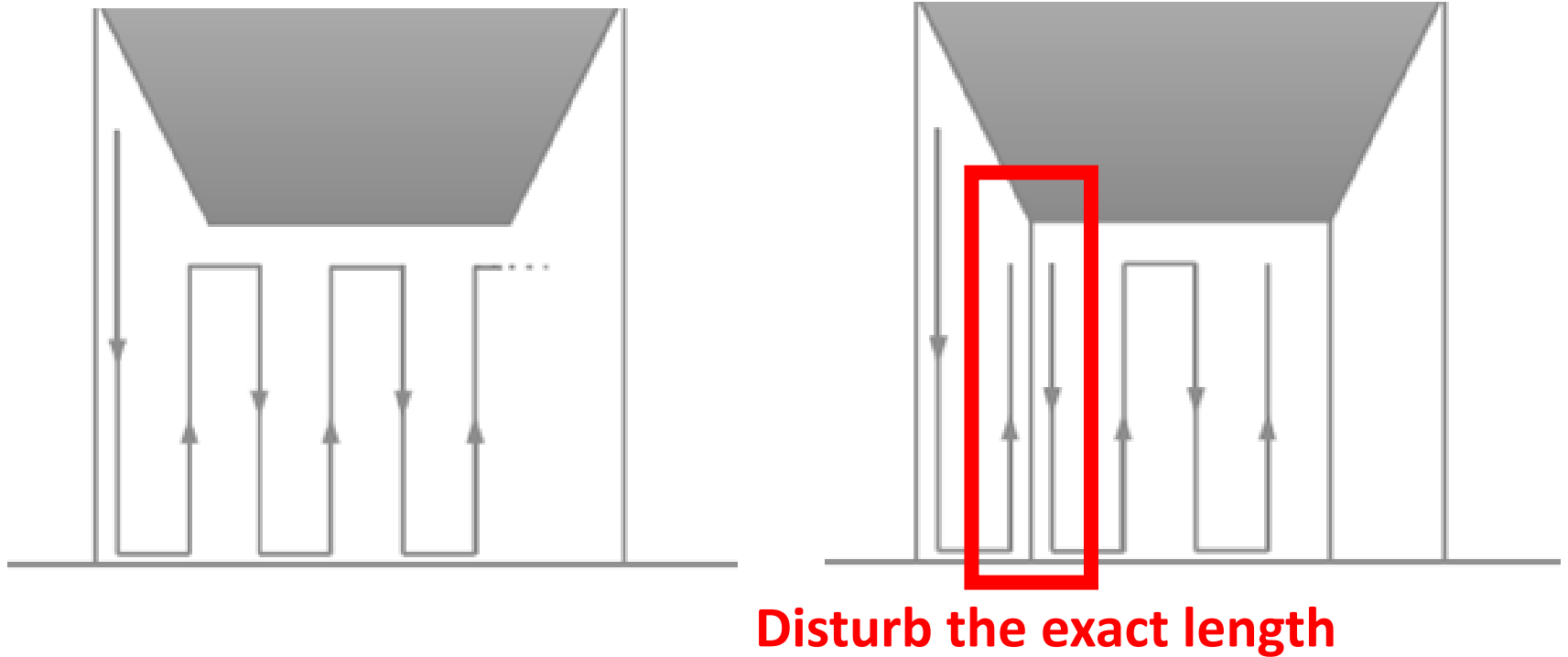


[Previous Work]

Boustrophedon Decomposition



Why Less Cells Are Better?



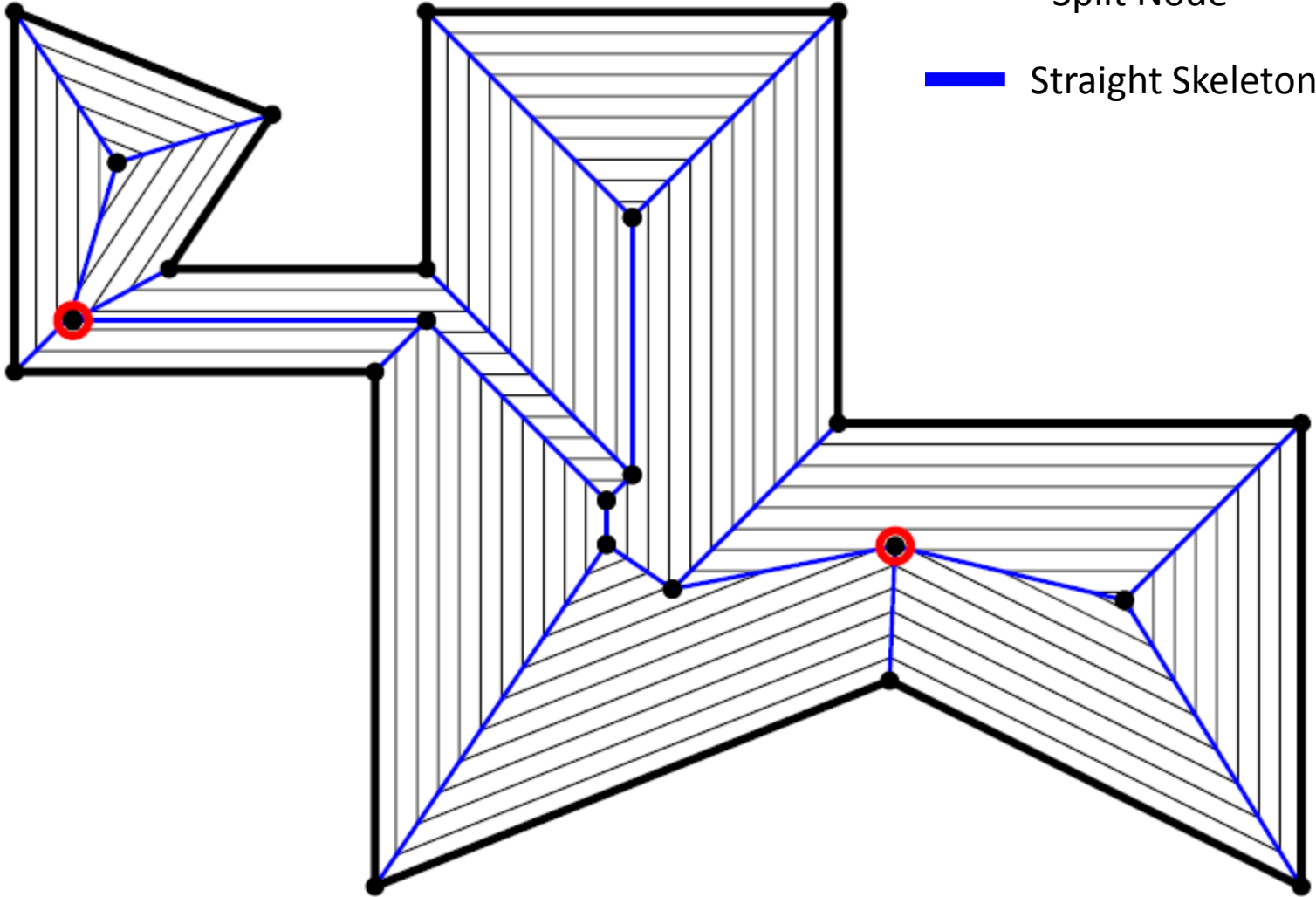
Constriction Decomposition Method

Constriction Decomposition Method

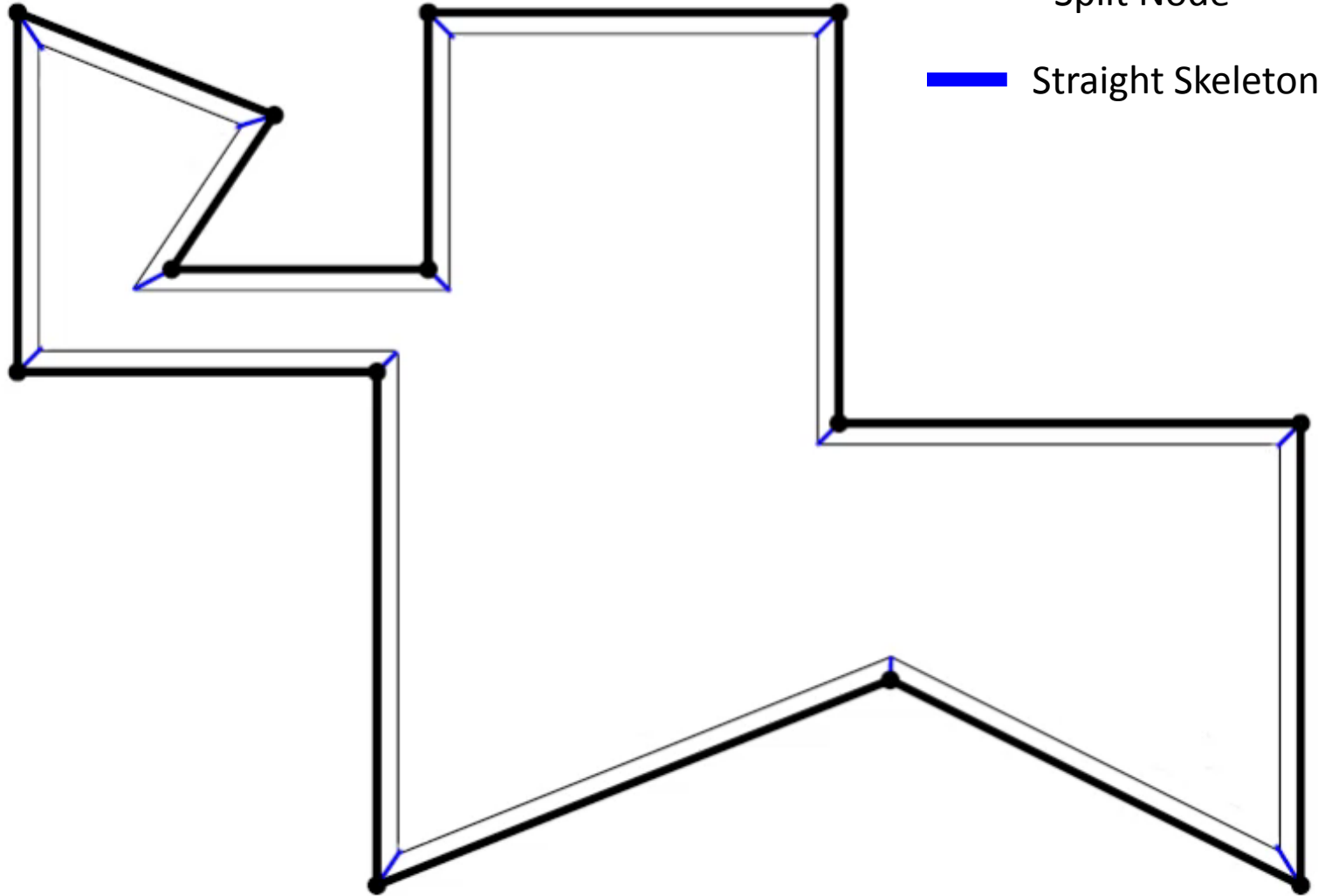
1. Generate a Straight Skeleton.
2. For all split nodes,
Generate a separator from a neighbor node of split nodes on boundary.

● Constriction Point
= Split Node

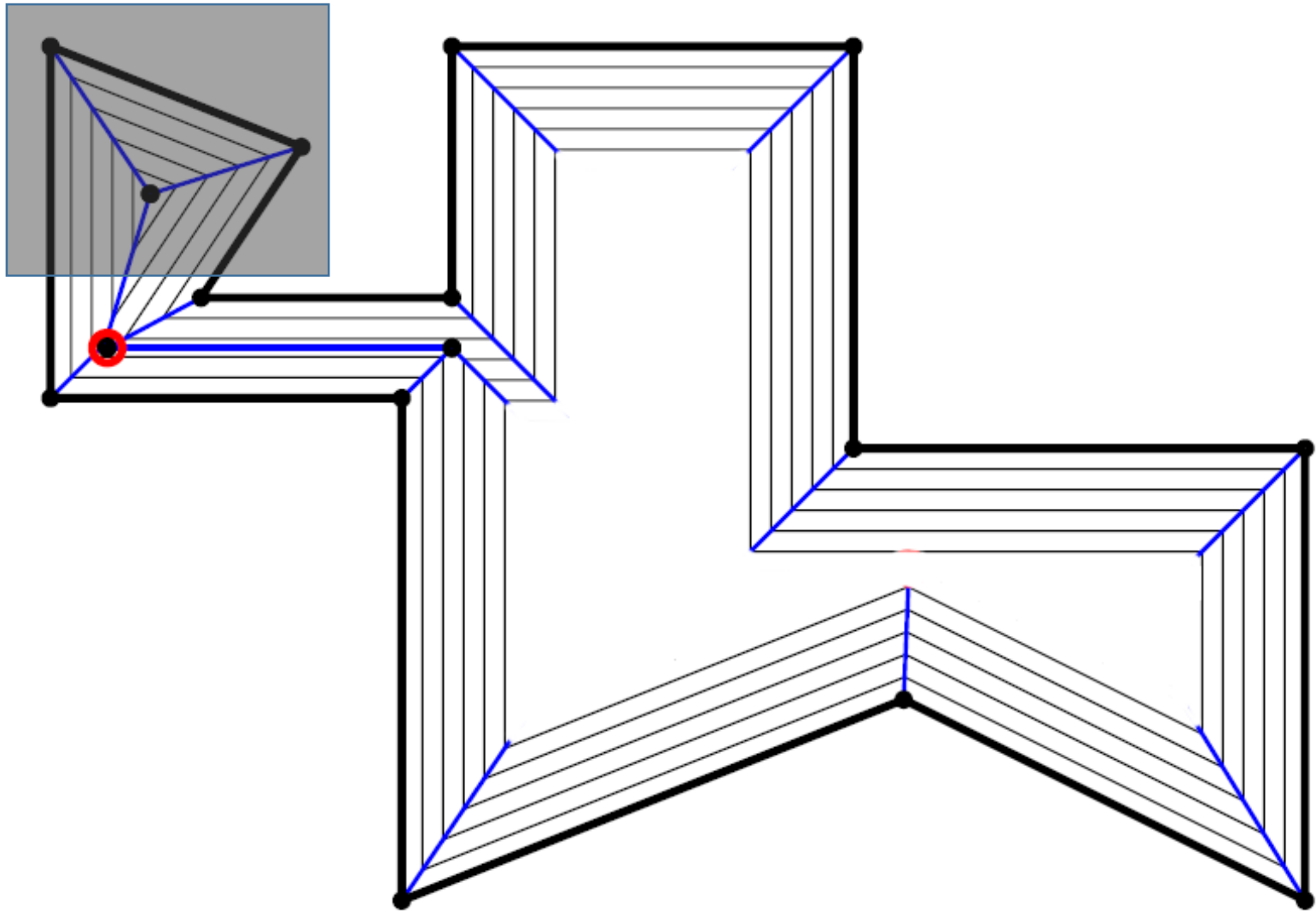
— Straight Skeleton

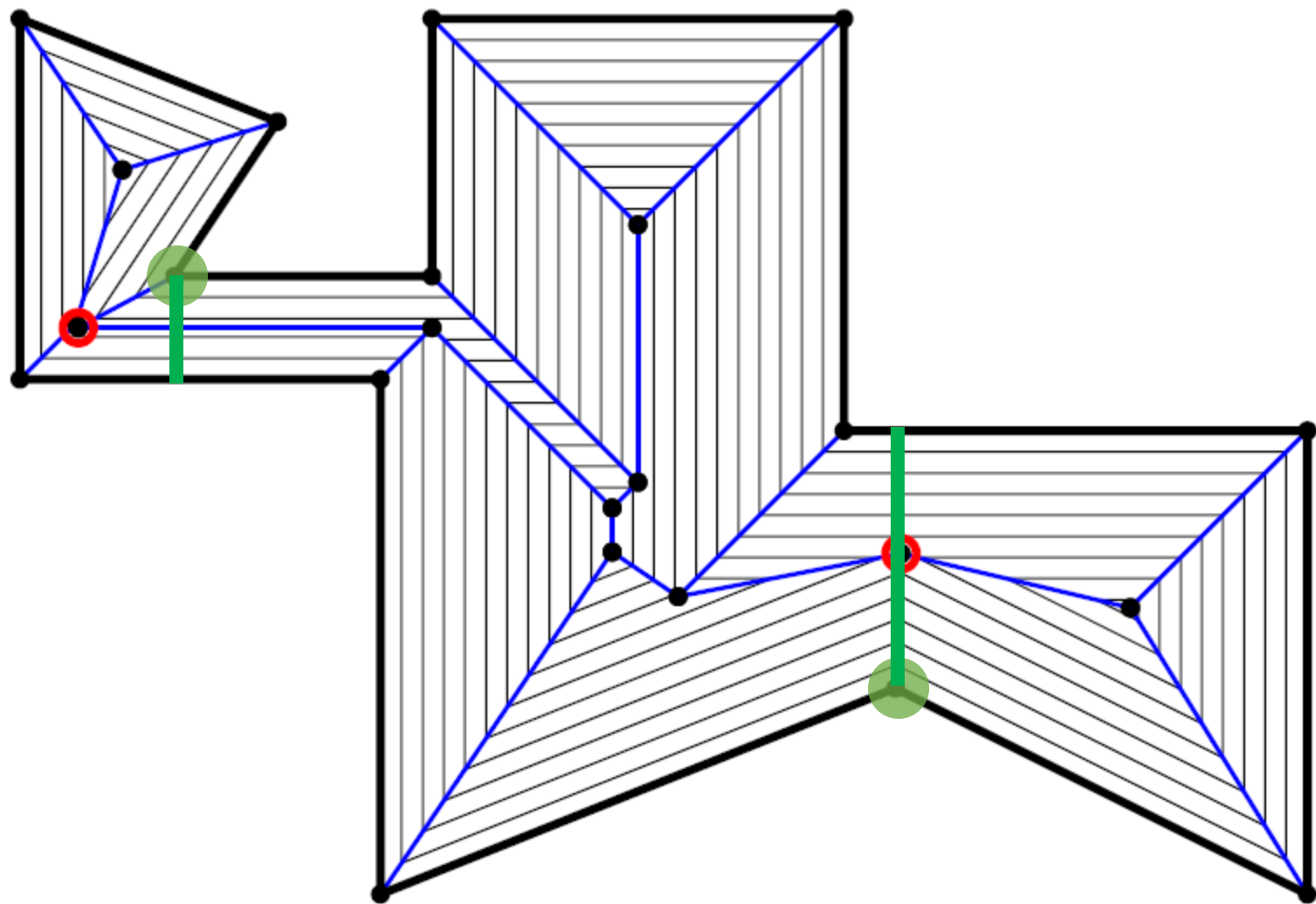


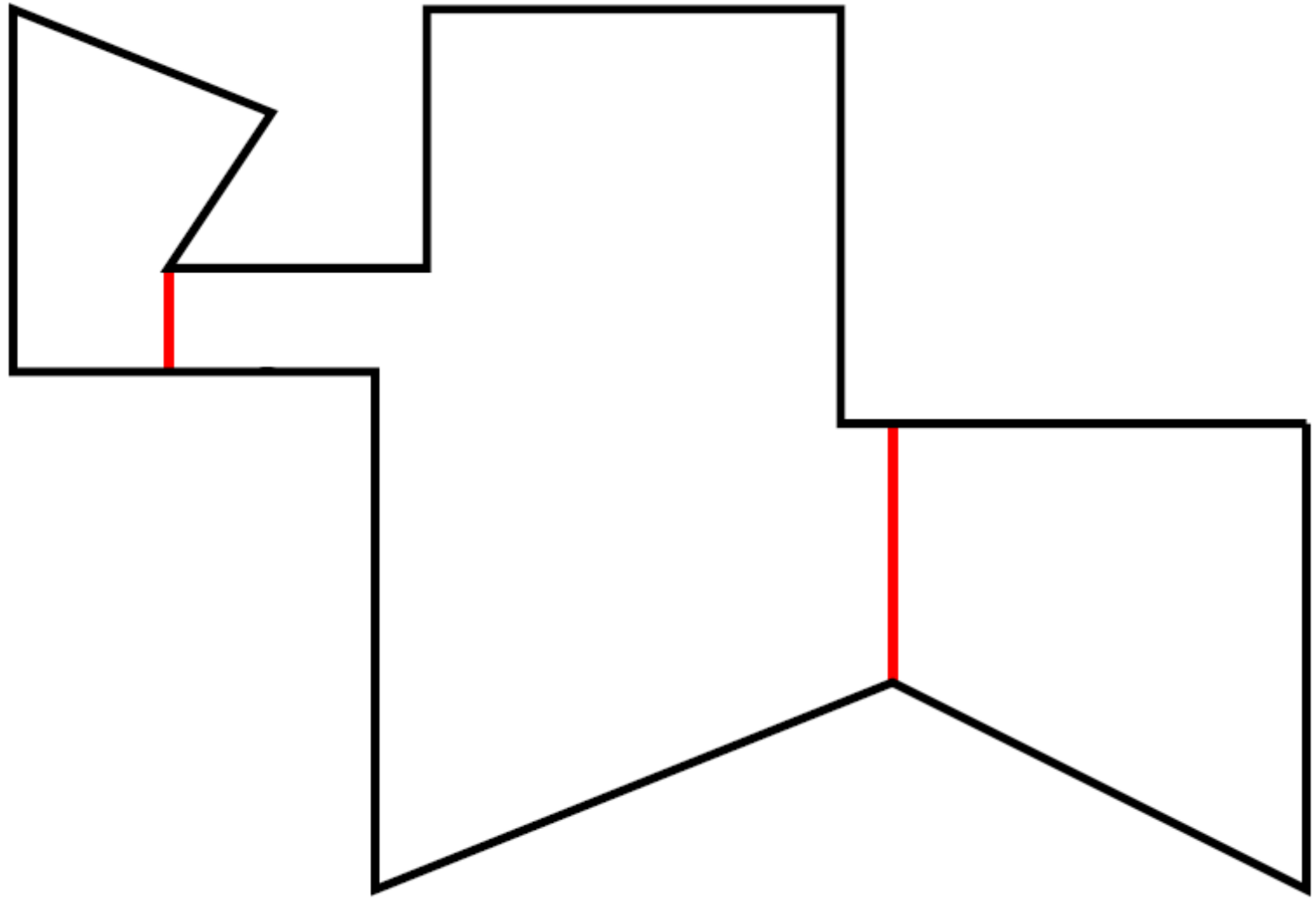
<Shrinking Process>



Edge Event

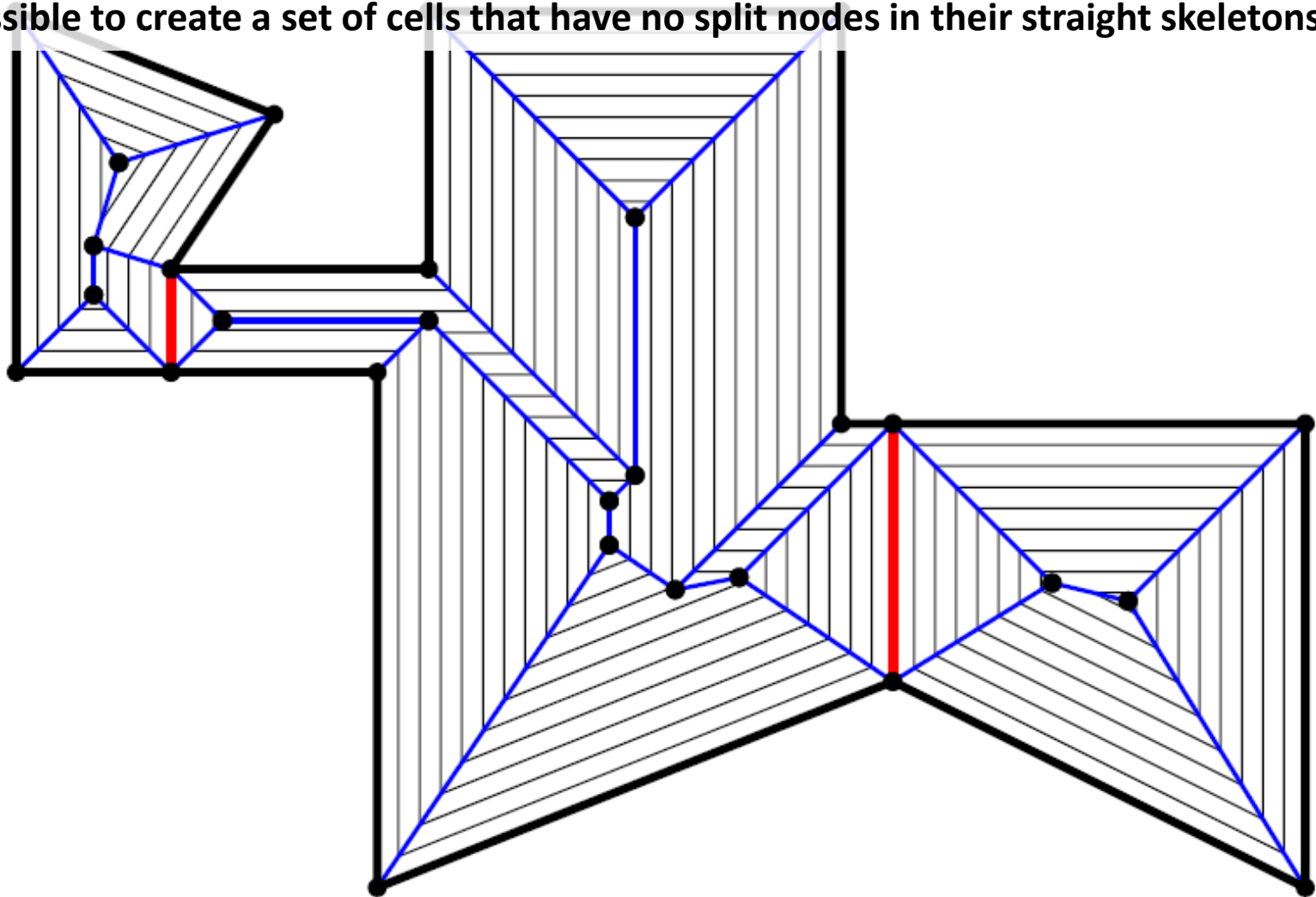




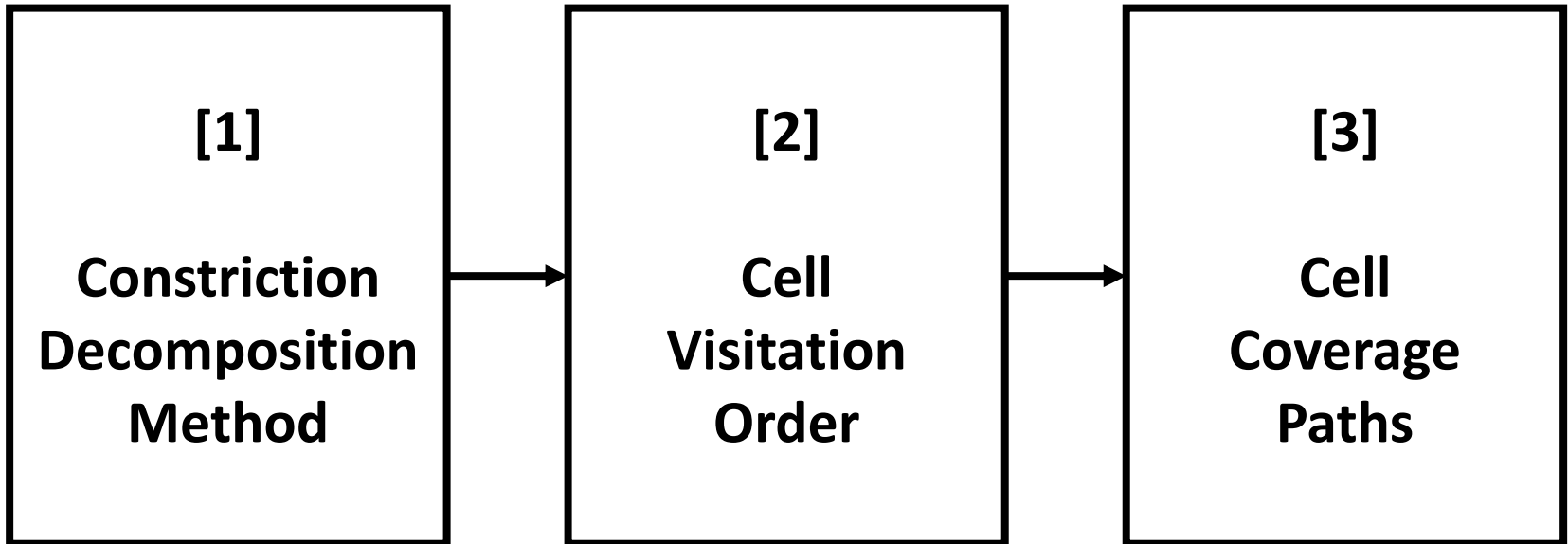


Can't Find Mathematical Background..

If the polygon is decomposed based on the split nodes in its straight skeleton, It is possible to create a set of cells that have no split nodes in their straight skeletons.

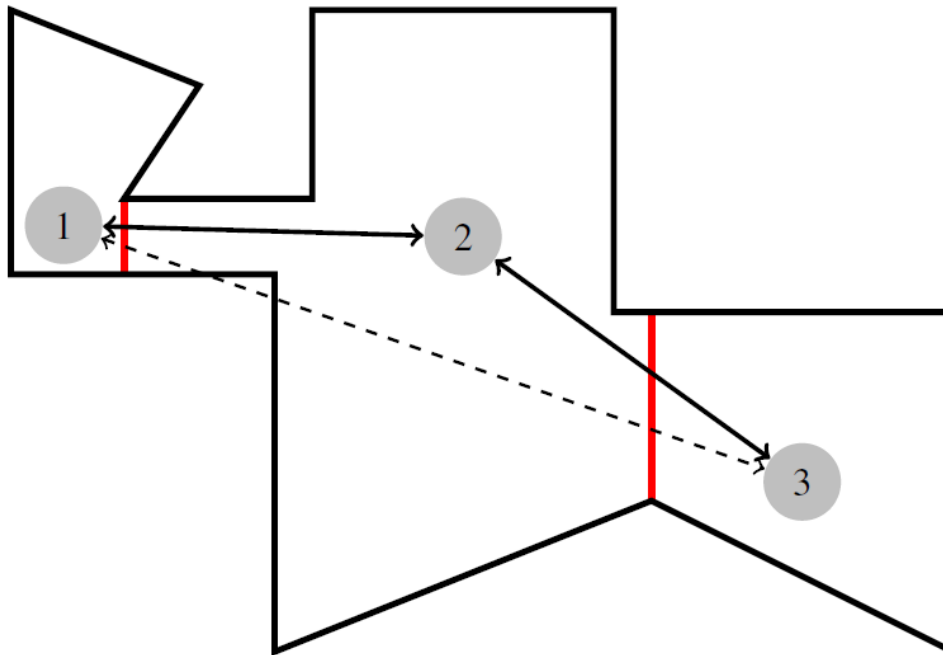


[Coverage Path Planning]



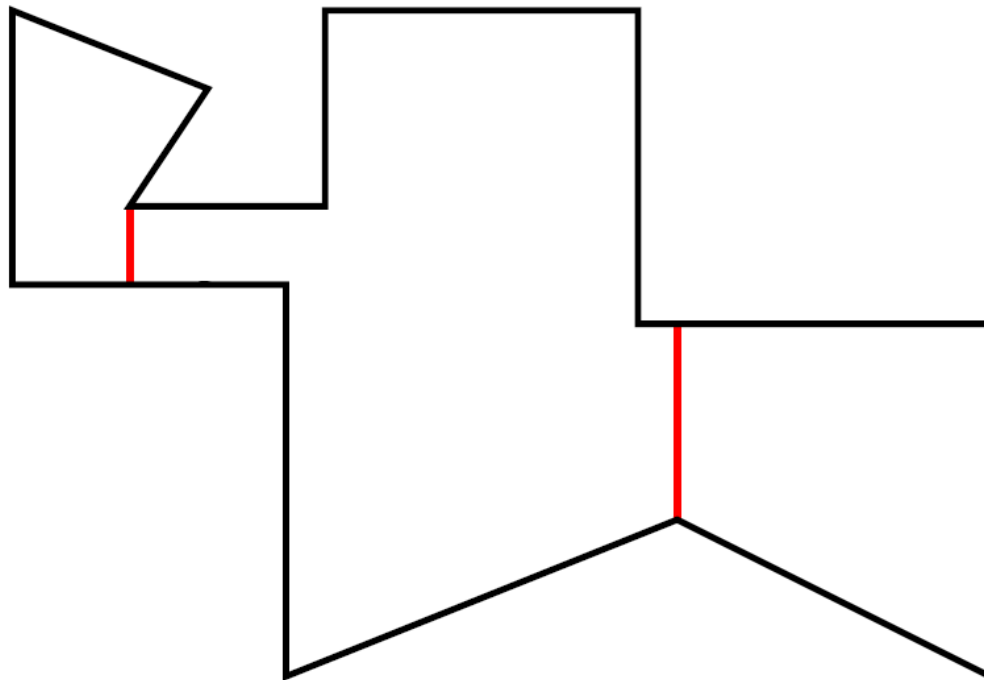
[2] Cell Visitation Order

- Should visit every cell.
- Generate an augmented adjacency graph.
- Solve using a *Heuristic TSP* algorithm.



[3] Cell Coverage Paths

- Applying Boustrophedon Decomposition to non-convex polygon directly might lead to over decomposition depending on the incline of the scan line.



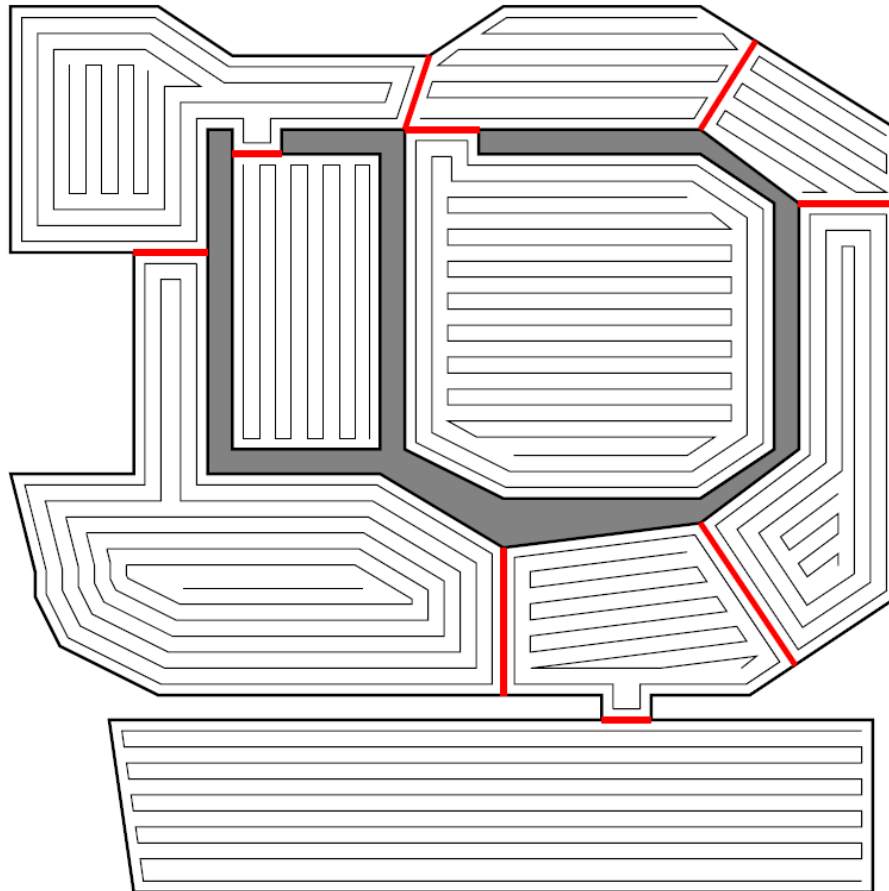
Results

Environment	BSD	CDM	Improvement
Non-convex	3	3	0%
Non-convex, with holes 1	14	9	36%
Non-convex, with holes 2	11	3	73%
Non-convex, with holes 3	9	5	44%
Floor Plan 1	169	53	68%
Floor Plan 2	158	65	58%
Floor Plan 3	172	66	62%

BSD: Boustrophedon Decomposition
CDM: Constriction Decomposition Method

Results

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Non-convex, with holes 1	14	9	36%



Results

Environment	BSD	CDM	Improvement
Floor Plan 1	169	53	68%

